

APPLICATION NOTE

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An Intelligent Data Base System Using the 8272

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1. INTRODUCTION

Most microcomputer systems in use today require lowcost, high-density removable magnetic media for information storage. In the area of removable media, a designer's choice is limited to magnetic tapes and floppy disks, (flexible diskettes), both of which offer non-volatile data storage. The choice between these two technologies is relatively straight-forward for a given application. Since disk drives are designed to permit random access to stored information, they are significantly faster than tape units. For example, locating information on a disk requires less than a second, while tape movement (even at the fastest rewind or fast-forward speed) often requires several minutes. This random access ability permits the use of floppy disks in on-line storage applications (where information must be located, read, and modified/updated in real-time under program or operator control). Tapes, on the other hand, are ideally suited to archival or back-up storage due to their large storage capacities (more than 10 million bytes of data can be archived on a cartridge tape).

A sophisticated controller is required to capitalize on the abilities of the disk storage unit. In the past, disk controller designs have required upwards of 150 ICs. Today, the single-chip 8272 Floppy Disk Controller (FDC) plus approximately 30 support devices can handle up to four million bytes of on-line data storage on four floppy disk drives.

The Floppy Disk

A floppy disk is a circular piece of thin plastic material covered with a magnetic coating and enclosed in a protective jacket (Figure 1). The circular piece of plastic revolves at a fixed speed (approximately 360 rpm) within its jacket in much the same manner that a record revolves at a fixed speed on a stereo turntable. Disks are manufactured in a variety of configurations for various storage capacities. Two standard physical disk sizes are commonly used. The 8-inch disk (8 inches square) is the larger of the two sizes; the smaller size (5-1/4 inches square) is often referred to as a mini-floppy. Singlesided disks can record information on only one side of the disk, while double-sided disks increase the storage capacity by recording on both sides. In addition, disks are classified as single-density or double-density. Doubledensity disks use a modified recording method to store twice as much information in the same disk area as can be stored on a single-density disk. Table 1 lists storage capacities for standard floppy disk media.

A magnetic head assembly (in contact with the disk) writes information onto the disk surface and subsequently reads the data back. This head assembly can move from the outside edge of the disk toward the center in fixed increments. Once the head assembly is

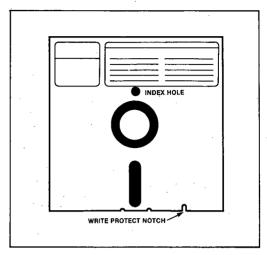


Figure 1. A Floppy Diskette

positioned at one of these fixed positions, the head can read or write information in a circular path as the disk revolves beneath the head assembly. This method divides the surface into a fixed number of cylinders (as shown in Figure 2). There are normally 77 cylinders on a standard disk. Once the head assembly is positioned at a given cylinder, data may be read or written on either side of the disk. The appropriate side of the disk is selected by the read/write head address (zero or one). Of course, a single-sided disk can only use head zero. The combination of cylinder address and head address uniquely specifies a single circular track on the disk. The physical beginning of a track is located by means of a small hole (physical index mark) punched through the plastic near the center of the disk. This hole is optically sensed by the drive on every revolution of the disk.

Table 1. Formatted Disk Capacities

Single-Density Format				,
Byte/Sector	128	256	512	1024
Sectors/Track	26	15	. 8	4
Tracks/Disk	77	. 77	77	77
Bytes/Disk	256,256	295,680	315,392	315,392
Double-Density Format				
Bytes/Sector	128	256	512	1024
Sectors/Track	52	30	16	- 8
Tracks/Disk	77	77	77	77
Bytes/Disk	512,512	591,360	630,784	630,784

Each track is subdivided into a number of sectors (see detailed discussion in section 3). Sectors are generally 128, 256, 512, or 1024 data bytes in length. This track sectoring may be accomplished by one of two techniques: hard sectoring or soft sectoring. Hard sectored disks divide each track into a maximum of 32 sectors. The beginning of each sector is indicated by a sector hole punched in the disk plastic. Soft sectoring, the IBM standard method, allows software selection of sector sizes. With this technique, each data sector is preceded by a unique sector identifier that is read/written by the disk controller.

A floppy disk may also contain a write protect notch punched at the edge of the outer jacket of the disk. This notch is detected by the drive and passed to the controller as a write protect signal.

The Floppy Disk Drive

The floppy disk drive is an electromechanical device that records data on, or reads data from, the surface of a floppy disk. The disk drive contains head control electronics that move the head assembly one increment (step) forward (toward the center of the disk) or backward (toward the edge of the disk). Since the recording head must be in contact with the disk material in order to read or write information, the disk drive also contains head-load electronics. Normally the read/write head is unloaded until it is necessary to read or write information on the floppy disk. Once the head assembly has been positioned over the correct track on the disk. the head is loaded (brought into contact with the disk). This sequence prevents excessive disk wear. A small time penalty is paid when the head is loaded. Approximately thirty to fifty milliseconds are needed before data may be reliably read from, or written to, the disk. This time is known as the head load time. If desired, the head may be moved from cylinder to cylinder while loaded. In this manner, only a small time interval (head settling time) is required before data may be read from the new cylinder. The head settling time is often shorter than the head load time. Typically, disk drives also contain drive select logic that allows more than one physical drive to be connected to the same interface cable (from the controller). By means of a jumper on the drive, the drive number may be selected by the OEM or end user. The drive is enabled only when selected; when not selected, all control signals on the cable are ignored.

Finally, the drive provides additional signals to the system controller regarding the status of the drive and disk. These signals include:

Drive Ready — Signals the system that the drive door is closed and that a floppy disk is inserted into the drive.

Track Zero — Indicates that the head assembly is located over the outermost track of the disk. This signal may be used for calibration of the disk drive at system initialization and after an error condition.

Write Protect — Indicates that the floppy disk loaded into the drive is write protected.

Dual Sided — Indicates that the floppy disk in the drive is dual-sided.

Write Fault — Indicates that an error occurred during a recording operation.

Index — Informs the system that the physical index mark of the floppy disk (signifying the start of a data track) has been sensed.

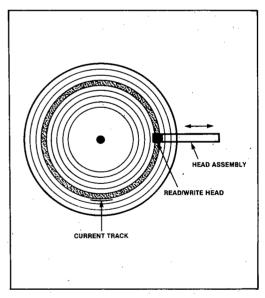


Figure 2. Concentric Cylinders on a Floppy Diskette

2. SUBSYSTEM OVERVIEW

A disk subsystem consists of the following functional electronic units:

- 1. Disk Controller Electronics
- 2. Disk Drive Electronics
- Controller/Disk Interface (cables, drivers, terminators)
- 4. Controller/Microprocessor System Interface

The operation of these functional units is discussed in the following paragraphs.

Controller Electronics

The disk controller is responsible for converting highlevel disk commands (normally issued by software executing on the system processor) into disk drive commands. This function includes:

- Disk Drive Selection Disk controllers typically manage the operations of multiple floppy disk drives. This controller function permits the system processor to specify which drive is to be used in a particular operation.
- 2. Track Selection The controller issues a timed sequence of step pulses to move the head from its current location to the proper disk cylinder from which data is to be read or to which data is to be written. The controller stores the current cylinder number and computes the stepping distance from the current cylinder to the specified cylinder. The controller also manages the head select signal to select the correct side of the floppy disk.
- Sector Selection The controller monitors the data on a track until the requested sector is sensed.
- 4. Head Loading The disk controller determines the times at which the head assembly is to be brought in contact with the disk surface in order to read or write data. The controller is also responsible for waiting until the head has settled before reading or writing information. Often the controller maintains the head loaded condition for up to 16 disk revolutions (approximately 2 seconds) after a read or write operation has been completed. This feature eliminates the head load time during periods of heavy disk I/O activity.
- 5. Data Separation The actual signal recorded on a floppy disk is a combination of timing information (clock) and data. The serial READ DATA input (from the disk drive) must be converted into two signal streams: clock and data. (The READ DATA input operates at 250K bits/second for single-density disks and 500K bits/second for double-density

- disks.) The serial data must also be assembled into 8-bit bytes for transfer to system memory. A byte must be assembled and transferred every 32 microseconds for single-density disks and every 16 microseconds for double-density.
- 6. Error Checking Information recorded on a floppy disk is subject to both hard and soft errors. Hard (permanent) errors are caused by media defects. Soft errors, on the other hand, are temporary errors caused by electromagnetic noise or mechanical interference. Disk controllers use a standard error checking technique known as a Cyclic Redundancy Check (CRC). As data is written to a disk, a 16-bit CRC character is computed and also stored on the disk. When the data is subsequently read, the CRC character allows the controller to detect data errors. Tvpically, when CRC errors are detected, the controlling software retries the failed operation (attempting to recover from a soft error). If data cannot reliably be read or written after a number of retries, the system software normally reports the error to the operator. Multiple CRC errors normally indicate unrecoverable media error on the current disk track. Subsequent recovery attempts must be defined by the system designers and tailored to meet system interfacing requirements.

Today, single-chip digital LSI floppy disk controllers such as the 8272 perform all the above functions with the exception of data separation. A data separation circuit (a combination of digital and analog electronics) synchronizes itself to the actual data rate of the disk drive. This data rate varies from drive to drive (due to mechanical factors such as motor tolerances) and varies from disk to disk (due to temperature effects). In order to operate reliably with both single- and double-density storage, the data separation circuit must be based on phase-locked loop (PLL) technology. The phase-locked loop data separation logic is described in section 5. The separation logic, after synchronizing with the data stream, supplies a data window to the LSI disk controller. This window differentiates data information from clock information within the serial stream. The controller uses this window to reconstruct the data previously recorded on the floppy disk.

Drive Electronics

Each floppy disk drive contains digital electronic circuits that translate TTL-compatible command signals into electromechanical operations (such as drive selection and head movement/loading) and that sense and report disk or drive status to the controller (e.g., drive ready, write fault, and write protect). In addition, the drive electronics contain analog components to sense, amplify, and shape data pulses read from, or written to, the floppy disk surface by the read/write head.

Controller/Drive Interface

The controller/drive interface consists of high-current line drivers, Schmitt triggered input gates, and flat or twisted pair cable(s) to connect the disk drive electronics to the controller electronics. Each interface signal line is resistively terminated at the end of the cable farthest from the line drivers. Eight-inch drives may be directly interfaced by means of 50-conductor flat cable. Generally, cable lengths should be less than ten feet in order to maintain noise immunity.

Normally, provisions are made for up to four disk drives to share the same interface cable. The controller may operate as many cable assemblies as practical. LSI floppy disk controllers typically operate one to four drives on a single cable.

Processor/Memory Interface

The disk controller must interface to the system processor and memory for two distinct purposes. First, the processor must specify disk control and command parameters to the controller. These parameters include the selection of the recording density and specification of disk formatting information (discussed in section 3). In addition to disk parameter specification, the processor must also send commands (e.g., read, write, seek, and scan) to the controller. These commands require the specification of the command code, drive number, cylinder address, sector address, and head address. Most LSI controllers receive commands and parameters by means of processor I/O instructions.

In addition to this I/O interface, the controller must also be designed for high-speed data transfer between memory and the disk drive. Two implementation methods may be used to coordinate this data transfer. The lowest-cost method requires direct processor intervention in the transfer. With this method, the controller issues an interrupt to the processor for each data transfer. (An equivalent method allows the processor to poll an interrupt flag in the controller status word.) In the case of a disk write operation, the processor writes a data byte (to be encoded into the serial output stream) to the disk controller following the receipt of each controller interrupt. During a disk read operation, the processor reads a data byte (previously assembled from the input data stream) from the controller after each interrupt. The processor must transfer a data byte from the controller to memory or transfer a data byte from memory to the disk controller within 16 or 32 microseconds after each interrupt (double-density and single-density response times, respectively).

If the system processor must service a variety of other interrupt sources, this interrupt method may not be practical, especially in double-density systems. In this case, the disk controller may be interfaced to a Direct Memory Access (DMA) controller. When the disk controller requires the transfer of a data byte, it simply activates the DMA request line. The DMA controller interfaces to the processor and, in response to the disk controller's request, gains control of the memory interface for a short period of time—long enough to transfer the requested data byte to/from memory. See section 6 for a detailed DMA interface description.

3. DISK FORMAT

New floppy disks must be written with a fixed format by the controller before these disks may be used to store data. Formatting is a method of taking raw media and adding the necessary information to permit the controller to read and write data without error. All formatting is performed by the disk controller on a track-bytrack basis under the direction of the system processor. Generally, a track may be formatted at any time. However, since formatting "initializes" a complete disk track, all previously written data is lost (after a format operation). A format operation is normally used only when initializing new floppy disks. Since soft-sectoring in such a predominant formatting technique (due to IBM's influence), the following discussion will limit itself to soft-sectored formats.

Data Recording Techniques

Two standard data recording techniques are used to combine clock and data information for storage on a floppy disk. The single-density technique is referred to as FM encoding. In FM encoding (see Figure 3), a double frequency encoding technique is used that inserts a data bit between two adjacent clock bits. (The presence of a data bit represents a binary "one" while the absence of a data bit represents a binary "zero.") The two adjacent clock bits are referred to as a bit cell, and except for unique field identifiers, all clock bits written on the disk are binary "ones." In FM encoding, each data bit is written at the center of the bit cell and the clock bits are written at the leading edge of the bit cell.

The encoding used for double-density recording is termed MFM encoding (for "Modified FM"). In MFM encoding (Figure 3) the data bits are again written at the center of the bit cell. However, a clock bit is written at the leading edge of the bit cell only if no data bit was written in the previous bit cell and no data bit will be written in the present bit cell.

Sectors

Soft-sectored floppy disks divide each track into a number of data sectors. Typically, sector sizes of 128, 256, 512, or 1024 data bytes are permitted. The sector size is specified when the track is initially formatted by the controller. Table 1 lists the single- and double-

density data storage capacities for each of the four sector sizes. Each sector within a track is composed of the following four fields (illustrated in Figure 4):

- 1. Sector ID Field This field, consisting of seven bytes, is written only when the track is formatted. The ID field provides the sector identification that is used by the controller when a sector must be read or written. The first byte of the field is the ID address mark, a unique coding that specifies the beginning of the ID field. The second, third, and fourth bytes are the cylinder, head, and sector addresses, respectively, and the fifth byte is the sector length code. The last two bytes are the 16-bit CRC character for the ID field. During formatting, the controller supplies the address mark. The cylinder, head, and sector addresses and the sector length code are supplied to the controller by the processor software. The CRC character is derived by the controller from the data in the first five bytes.
- 2. Post ID Field Gap The post ID field gap (gap 2) is written initially when the track is formatted. During subsequent write operations, the drive's write circuitry is enabled within the gap and the trailing bytes of the gap are rewritten each time the sector is updated (written). During subsequent read operations, the trailing bytes of the gap are used to synchronize the data separator logic with the upcoming data field.
- 3. Data Field The length (number of data bytes) of the data field is determined by software when the track is formatted. The first byte of the data field is the data address mark, a unique coding that specifies

- the beginning of the data field. When a sector is to be deleted, (e.g., a hard error on the disk), a deleted data address mark is written in place of the data address mark. The last two bytes of the data field comprise the CRC character.
- 4. Post Data Field Gap The post data field gap (gap 3) is written when the track is formatted and separates the preceding data field from the next physical ID field on the track. Note that a post data field gap is not written following the last physical sector on a track. The gap itself contains a programselectable number of bytes. Following a sector update (write) operation, the drive's write logic is disabled during the gap. The actual size of gap 3 is determined by the maximum number of data bits that can be recorded on a track, the number of sectors per track and the total sector size (data plus overhead information). The gap size must be adjusted so that it is large enough to contain the discontinuity generated on the floppy disk when the write current is turned on or off (at the start or completion of a disk write operation) and to contain a synchronization field for the upcoming ID field (of the next sector). On the other hand, the gaps must be small enough so that the total number of data bits required on the track (sectors plus gaps) is less than the maximum number of data bits that can be recorded on the track. The gap size must be specified for all read, write, and format operations. The gap size used during disk reads and writes must be smaller than the size used to format the disk to avoid the splice points between contiguous physical sectors. Suggested gap sizes are listed in Table 9.

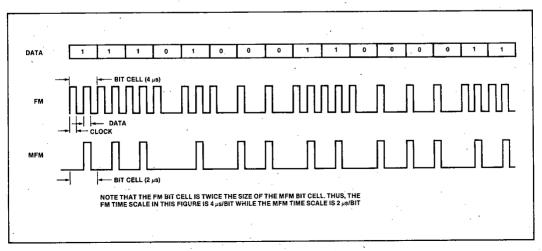


Figure 3. FM and MFM Encoding

Tracks

The overall format for a track is illustrated in Figure 4. Each track consists of the following fields:

- Pre-Index Gap The pre-index gap (gap 5) is written only when the track is formatted.
- Index Address Mark The index address mark consists of a unique code that indicates the beginning of a data track. One index mark is written on each track when the track is formatted.
- Post Index Gap The post index gap (gap 1) is used during disk read and write operations to syn-

- chronize the data separator logic with the data to be read from the ID field (of the first sector). The post index gap is written only when the disk is formatted.
- 4. Sectors The sector information (discussed above) is repeated once for each sector on the track.
- 5. Final Gap The final gap (gap 4) is written when the track is formatted and extends from the last physical data field on the track to the physical index mark. The length of this gap is dependent on the number of bytes per sector specified, the lengths of the program-selectable gaps specified, and the drive speed.

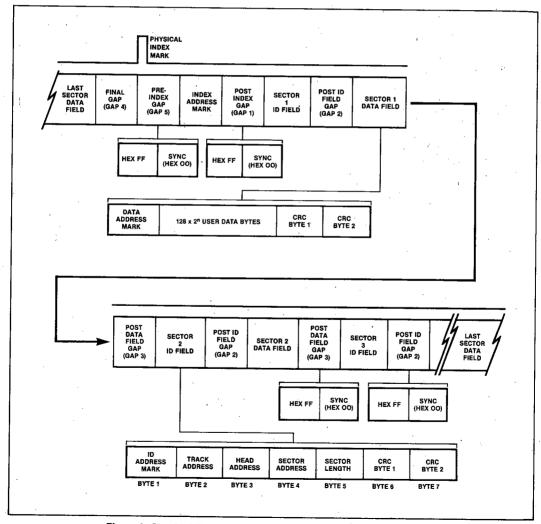


Figure 4. Standard Floppy Diskette Track Format (From SBC 204 Manual)

Sector Interleaving

The initial formatting of a floppy disk determines where sectors are located within a track. It is not necessary to allocate sectors sequentially around the track (i.e., 1,2,3,...,26). In fact, is is often advantageous to place the sectors on the track in a non-sequential order. Sequential sector ordering optimizes sector access times during multi-sector transfers (e.g., when a program is loaded) by permitting the number of sectors specified (up to an entire track) to be transferred within a single revolution of the disk. A technique known as sector interleaving optimizes access times when, although sectors are accessed sequentially, a small amount of processing must be performed between sector reads/writes. For example, an editing program performing a text search reads sectors sequentially, and after each sector is read, performs a software search. If a match is not found, the software issues a read request for the next sector. Since the floppy disk continues to rotate during the time that the software executes, the next physical sector is already passing under the read/write head when the read request is issued, and the processor must wait for another complete revolution of the disk (approximately 166 milliseconds) before the data may actually be input. With interleaving, the sectors are not stored sequentially on a track; rather, each sector is physically removed from the previous sector by some number (known as the interleave factor) of physical sectors as shown in Figure 5. This method of sector allocation provides the processor additional execution time between sectors on the disk. For example, with a 26 sector/track format, an interleave factor of 2 provides 6.4 milliseconds of processing time between sequential 128 byte sector accesses.

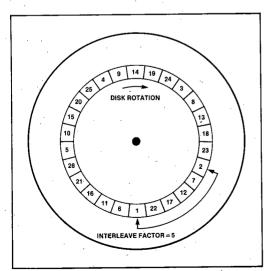


Figure 5. Interleaved Sector Allocation Within a Track

To calculate the correct interleave factor, the maximum processor time between sector operations must be divided by the time required for a complete sector to pass under the disk read/write head. After determining the interleave factor, the correct sector numbers are passed to the disk controller (in the exact order that they are to physically appear on the track) during the execution of a format operation.

4. THE 8272 FLEXIBLE DISKETTE CONTROLLER

The 8272 is a single-chip LSI Floppy Disk Controller (FDC) that contains the circuitry necessary to implement both single-and double-density floppy disk storage subsystems (with up to four dual-sided disk drives per FDC). The 8272 supports the IBM 3740 single-density recording format (FM) and the IBM System 34 doubledensity recording format (MFM). With the 8272, less than 30 ICs are needed to implement a complete disk subsystem. The 8272 accepts and executes high-level disk commands such as format track, seek, read sector, write sector, and read track. All data synchronization and error checking is automatically performed by the FDC to ensure reliable data storage and subsequent retrieval. External logic is required only for the generation of the FDC master clock and write clock (see Section 6) and for data separation (Section 5). The FDC provides signals that control the startup and base frequency selection of the data separator. These signals greatly ease the design of a phase-locked loop data separator.

In addition to the data separator interface signals, the 8272 also provides the necessary signals to interface to microprocessor systems with or without Direct Memory Access (DMA) capabilities. In order to interface to a large number of commercially available floppy disk drives, the FDC permits software specification of the track stepping rate, the head load time, and the head unload time.

The pin configuration and internal block diagram of the 8272 is shown in Figure 6. Table 2 contains a description for each FDC interface pin.

Floppy Disk Commands

The 8272 executes fifteen high-level disk interface commands:

Specify
Sense Drive Status
Sense Interrupt Status
Seek
Recalibrate
Format Track
Read Data
Read Deleted Data

Write Data Write Deleted Data Read Track Read ID Scan Equal

Scan High or Equal Scan Low or Equal

Each command is initiated by a multi-byte transfer from the processor to the FDC (the transferred bytes contain command and parameter information). After complete command specification, the FDC automatically executes the command. The command result data (after execution of the command) may require a multi-byte transfer of status information back to the processor. It is convenient to consider each FDC command as consisting of the following three phases:

COMMAND PHASE:

The executing program transfers to the FDC all the information required to perform a particular disk operation. The 8272 automatically enters the command phase after RESET and following the completion of the result phase (if any) of a previous command.

EXECUTION PHASE:

The FDC performs the operation as instructed. The execution phase is entered immediately after the last command parameter is written to the FDC in the preceding command phase. The execution phase normally ends when the last data byte is transferred to/from the disk (signalled by the TC input to the FDC) or when an error occurs.

RESULT PHASE:

After completion of the disk operation, status and other housekeeping information are made available to the processor. After the processor reads this information, the FDC reenters the command phase and is ready to accept another command.

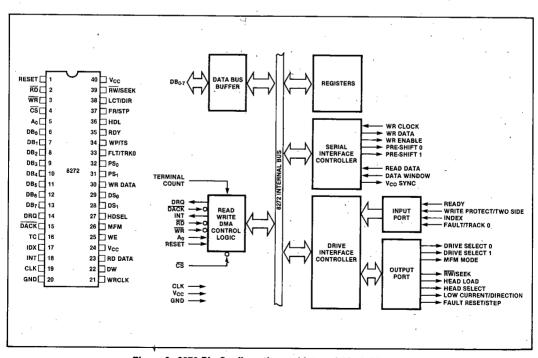


Figure 6. 8272 Pin Configuration and Internal Block Diagram

Table 2. 8272 FDC Pin Description

Number	Pin Symbol	I/O	To/From	Description
1	RST	I	uP	Reset. Active-high signal that places the FDC in the "idle" state and all disk drive output signals are forced inactive (low). This input must be held active during power on reset while the RD and WR inputs are active.
2	RD	Į*	uР	Read. Active-low control signal that enables data transfer from the FDC to the data bus.
3	WR	I*	uP	Write. Active-low control signal that enables data transfer from the data bus into the FDC.
, . 4	· CS	I	uP	Chip Select. Active-low control signal that selects the FDC. No reading or writing will occur unless the FDC is selected.
5	A ₀	I*	uP .	Address. Selects the Data Register or Main Status Register for input/out- put in conjunction with the RD and WR inputs. (See Table 3.)
6-13	DB ₀ -DB ₇	I/O*	uP	Data Bus. Bidirectional three-state 8-bit data bus.
14	DRQ	0	DMA	DMA Request. Active-high output that indicates an FDC request for DMA services.
15	DACK	I	DMA	DMA Acknowledge. Active-low control signal indicating that the requested DMA transfer is in progress.
16	TC	1	DMA	Terminal Count. Active-high signal that causes the termination of a command. Normally, the terminal count input is directly connected to the TC/EOP output from the DMA controller, signalling that the DMA transfer has been completed. In a non-DMA environment, the processor must count data transfers and supply a TC signal to the FDC.
17	IDX	I	Drive	Index. Indicates detection of the physical index mark (the beginning of a track) on the selected disk drive.
18	INT	0,	uP	Interrupt Request. Active-high signal indicating an 8272 interrupt service request.
19	CLK	I		Clock. Signal phase 8 MHz clock (50% duty cycle).
20	GND			Ground, DC power return.
21	WR CLK	I		Write Clock. 500 kHz (FM) or 1 MHz (MFM) write clock with a constant pulse width of 250 ns (for both FM and MFM recording). The write clock must be present at all times.
22	DW	I	PLL	Data Window. Data sample signal from the phase-locked loop indicating that the FDC should sample input data from the disk drive.
23	RD DATA	I.	Drive	Read Data. FDC input data from the selected disk drive.
24	VCO	0	PLL	VCO Sync. Active-high output that enables the phase-locked loop to synchronize with the input data from the disk drive.
25	WE	0	Drive	Write Enable. Active-high output that enables the disk drive write gate.
26	MFM	О	PLL	MFM Mode. Active-high output used by external logic to enable the MFM double-density recording mode. When the MFM output is low, single-density FM recording is indicated.
27	HDSEL	0.	Drive	Head Select. Selects head 0 or head 1 on a dual-sided disk.
28,29	DS ₁ ,DS ₀	o ·	Drive	Drive Select. Selects one of four disk drives.
30	WR DATA	0	Drive	Write Data. Serial data stream (combination of clock and data bits) to be written on the disk.
31,32	PS ₁ ,PS ₀	0	Drive	Precompensation (pre-shift) Control. Write precompensation output control during MFM mode. Specifies early, late, and normal timing signals. See the discussion in Section 5.

Table 2. 8272 FDC Pin Description (continued)

Number	Pin Symbol	1/0	To/From	Description				
33	FLT/TRKO	I	Drive	Fault/Track 0. Senses the disk drive fault condition in the Read/Write mode and the Track 0 condition in the Seek mode.				
. 34	WP/TS	I	Drive	Write Protect/Two-Sided. Senses the disk write protect status in the Read/Write mode and the dual-sided media status in the Seek mode.				
35	RDY	I	Drive	Ready. Senses the disk drive ready status.				
. 36	HDL	0	Drive	Head Load. Loads the disk drive read/write head. (The head is placed in contact with the disk.)				
37	FR/STP	0	Drive	Fault Reset/Step. Resets the fault flip-flop in the disk drive when operating in the Read/Write mode. Provides head step pulses (to move the head from one cylinder to another cylinder) in the Seek mode.				
38	LCT/DIR	0	Drive	Low Current/Direction. Signals that the recording head has been positioned over the inner cylinders (44-77) of the floppy disk in the Read/Write mode. (The write current must be lowered when recording on the physically shorter inner cylinders of the disk. Most drives do not track the actual head position and require that the FDC supply this signal.) Determines the head step direction in the Seek mode. In the Seek mode, a high level on this pin steps the read/write head toward the spindle (step-in); a low level steps the head away from the spindle (step-out).				
39	RW/SEEK	0	Drive	Read, Write/Seek Mode Selector. A high level selects the Seek mode; a low level selects the Read/Write mode.				
40	v_{cc}			+ 5V DC Power.				

^{*}Disabled when CS is high.

Interface Registers

To support information transfer between the FDC and the system processor, the 8272 contains two 8-bit registers: the Main Status Register and the Data Register. The Main Status Register (read only) contains FDC status information and may be accessed at any time. The Main Status Register (Table 4) provides the system processor with the status of each disk drive, the status of the FDC, and the status of the processor interface. The Data Register (read/write) stores data, commands, parameters, and disk drive status information. The Data Register is used to program the FDC during the command phase and to obtain result information after completion of FDC operations. Data is read from, or written to, the FDC registers by the combination of the A0, RD, WR, and CS signals, as described in Table 3.

In addition to the Main Status Register, the FDC contains four additional status registers (ST0, ST1, ST2, and ST3). These registers are only available during the result phase of a command.

Table 3. FDC Read/Write Interface

CS `	A ₀	RD	WR	Function
0	0	0	1	Read Main Status Register
0	0	1	0	Illegal
0	0	0	0	Illegal
0	1	0	0	Illegal
0	1	0	1	Read from Data Register
0	1	1	0	Write into Data Register
1	Х	X	X	Data Bus is three-stated

Table 4. Main Status Register Bit Definitions

Bit Number	Symbol	Description					
0	D ₀ B	Disk Drive 0 Busy. Disk Drive 0 is in the Seek mode.					
1	D ₁ B	Disk Drive 1 Busy. Disk Drive 1 is in the Seek mode.					
2	D ₂ B	Disk Drive 2 Busy. Disk Drive 2 is in the Seek mode.					
3	D ₃ B	Disk Drive 3 Busy. Disk Drive 3 is in the Seek mode.					
4	СВ	FDC Busy. A read or write command is in process.					
5	NDM	Non-DMA Mode. The FDC is in the non-DMA mode when this bit is high. This bit is set only during the execution phase of commands in the non-DMA mode. Transition to a low level indicates that the exe- cution phase has ended.					
6	DIO	Data Input/Output. Indicates the direction of a data transfer between the FDC and the Data Register. When DIO is high, data is read from the Data Register by the processor; when DIO is low, data is written from the processor to the Data Register.					
7	RQM	Request for Master. Indicates that the Data Register is ready to send data to, or receive data from, the processor.					

Command/Result Phases

Table 5 lists the 8272 command set. For each of the fifteen commands, command and result phase data transfers are listed. A list of abbreviations used in the table is given in Table 6, and the contents of the result status registers (ST0-ST3) are illustrated in Table 7.

The bytes of data which are sent to the 8272 during the command phase, and are read out of the 8272 in the result phase, must occur in the order shown in Table 5. That is, the command code must be sent first and the other bytes sent in the prescribed sequence. All bytes of the command and result phases must be read/written as described. After the last byte of data in the command phase is sent to the 8272 the execution phase automatically starts. In a similar fashion, when the last byte of data is read from the 8272 in the result phase,

the command is automatically ended and the 8272 is ready for a new command. A command may be aborted by simply raising the terminal count signal (pin 16). This is a convenient means of ensuring that the processor may always gain control of the 8272 (even if the disk system hangs up in an abnormal manner).

It is important to note that during the result phase all bytes shown in Table 5 must be read. The Read Data command, for example, has seven bytes of data in the result phase. All seven bytes must be read in order to successfully complete the Read Data command. The 8272 will not accept a new command until all seven bytes have been read. The number of command and result bytes varies from command-to-command.

In order to read data from, or write data to, the Data Register during the command and result phases, the system processor must examine the Main Status Register to determine if the Data Register is available. The DIO (bit 6) and RQM (bit 7) flags in the Main Status Register must be low and high, respectively, before each byte of the command word may be written into the 8272. Many of the commands require multiple bytes, and as a result, the Main Status Register must be read prior to each byte transfer to the 8272. To read status bytes during the result phase, DIO and RQM in the Main Status Register must both be high. Note, checking the Main Status Register in this manner before each byte transfer to/from the 8272 is required only in the command and result phases, and is NOT required during the execution phase.

Execution Phase

All data transfers to (or from) the floppy drive occur during the execution phase. The 8272 has two primary modes of operation for data transfers (selected by the specify command):

- 1. DMA mode
- 2. non-DMA mode

In the DMA mode, DRQ (DMA Request) is activated for each transfer request. The DMA controller responds to DRQ with \overline{DACK} (DMA Acknowledge) and \overline{RD} (for read commands) or \overline{WR} (for write commands). DRQ is reset by the FDC during the transfer. INT is activated after the last data transfer, indicating the completion of the execution phase, and the beginning of the result phase. In the DMA mode, the terminal count (TC/EOP) output of the DMA controller should be connected to the 8272 TC input to properly terminate disk data transfer commands.

Table 5. 8272 Command Set

		DATA BUS				DATA BUS	
PHASE	R/W	D ₇ D ₆ D ₅ D ₄ D ₃ D ₂ D ₁ D ₀	REMARKS .	PHASE	R/W	D ₇ D ₆ D ₅ D ₄ D ₃ D ₂ D ₁ D ₀	REMARKS
		READ DATA				READ A TRACK	
.Command	w	MT MFM SK 0 0 1 1 0	Command Codes	Command	w	0 MFM SK 0 0 0 1 0	Command Codes
	w	0 0 0 0 0 HDS DS1 DS0			w	0 0 0 0 0 HDS DS1 DS0	
- 1	W	С	Sector ID information prior to Command		w	C	Sector ID information prior to Command
	w	R	execution		. w	l R'	execution
	w	N EOT			w	N EOT	
	w	GPL			w	GPL	
	w	DTL			W	DTL	
Execution			Data transfer	Execution			Data transfer between the FDD
			between the FDD and the main-system			•	and the main-system.
Result	R	ST 0	Status information				FDC reads the complete track
	R	ST 1	after Command				contents from the
1	R	ST 2	execution				physical index mark to EOT
}	R	CH	Sector ID information		_		
i	R	R	after command execution	Result	R	ST 0 ST 1 ST 2	Status information after Command
لــنــا	_ n		execution		R	ST 2	execution
		READ DELETED DATA			A	<u>C</u>	Sector ID information
Command	W	MT MFM SK 0 1 1 0 0	Command Codes		R	H	after Command
	w	0 0 0 0 0 HDS DS1 DS0	Sector ID information		R	N	execution
	w.	Н	prior to Command			READ ID	
	w	Н	execution	Command	w	0 MFM 0 0 1 0 1 0	Command Codes
	W	N			w	0 0 0 0 0 HDS DS1 DS0	
	w	GPL					
	w	DIL		Execution		· ·	The first correct ID information on the
Execution			Data transfer between the FDD				track is stored in
			and the main-system			·	Data Register
Result	R.	STO	Status information	Result	R	ST 0	Status information
	Ř	ST 1 ST 2 C	after Command	[R R	ST 1 ST 2	after Command execution
	R	S12	execution	1	R	C	execution
	R	H	Sector ID information	1	R	н	Sector ID information
	R	R	after Command execution		R	R	during Execution Phase
		WRITE DATA	excedion			FORMAT A TRACK	
Command	w	MT MFM 0 0 0 1 0 1	Command Codes	Command	w	0 MFM 0 "0 1 1 0 1	Command Codes
Continand	w	0 0 0 0 0 HDS DS1 DS0	Command Codes	001111111111111111111111111111111111111	w	0 0 0 0 0 HDS DS1 DS0	Communa Codes
	w	c	Sector ID information		w	N	Bytes/Sector
	W	H	prior to Command	1 .	w	SC GPL	Sectors/Track Gap 3
	w		execution	1	w	D	Filter Byte
	w	EOT		Execution			#DO 11-
- 1	W	GPL DTL		Execution			FDC formats an entire track
Execution			Data transfer	l	_		
LXCCUITOIT			between the main-	Result	R	ST 0 ST 1 ST 2	Status information after Command
		i i	system and the FDD		R	ST 2	execution
Result	R	ST 0	Status information		R	C	In this case, the ID
	R	ST2	after Command execution		R	R	information has no
	R		i. I		R	N	meaning
	R R	H	Sector ID information after Command			SCAN EQUAL .	
	R	N	execution	Command	w	MT MFM SK 1 0 0 0 1	Command Codes
		WRITE DELETED DATA			w	0 0 0 0 0 HDS DS1 DS0	
Command	W	MT MFM 0 0 1 0 0 1	Command Codes		w	С	Sector ID information prior to Command
]	W	0 0 0 0 0 HDS DS1 DS0			w	R	execution
Ì	W		Sector ID information		w	NEOT	·
	w	H	prior to Command execution	'	w	GPL	
	W	N			w	STP	
	w	EOT GPL		Execution			Data compared
	w	DTL	,			·	between the FDD
		·	Data transfer				and the main system
Execution			between the FDD	Result	R	ST0	Status information
Execution			and the main-system		R	ST 1	after Command 'execution
					n 1		
Execution Result	R	ST 0	Status information	· [R	c	
	· A R	ST 1 ST 2	Status information after Command execution		R]	C	Sector ID information
	R R	ST 1 ST 2 C C	after Command execution		R	C H R	Sector ID information after Command
	· A R	ST 1 ST 2	after Command		R]	C	Sector ID information

Note: 1. A₀ = 1 for all operations.

Table 5. Command Set (Continued)

г	DATA BUS								1													
PHASE	R/W	D ₇	D ₆	D ₅	`D4	D ₃	D ₂	D	1 1	D _O	REMARKS	PHASE	R/W	07	D ₆	D ₅	D ₄	D ₃	D ₂	D ₁	D _O	REMARKS
				sc	AN L	ow c) R E	AUC	L								REC	ALIB	RATE			
Command	w	MT 0	•	0	0		HD	S DS	31 C	080	Command Codes	Command	w	0	0	0	0	0			1 DS0	Command Codes
	w					н					Sector ID information prior Command execution	Execution										Head retracted to Track 0
	w.					Ν					CACCUION				5	EN\$	E INT	ERRI	JPT S	TAT	JS	· · · · · · · · · · · · · · · · · · ·
	W				_ `E	OT _	_			_		Command	w						0			Command Codes
	w	=			s	TP _				_		Result	R	_			_ s	T0 -				Status information at the end of each seek operation about the
Execution					•					·	Data compared between the FDD and the main-system			<u>L</u> _								FDC
												٠, ,					S	PECI	FY			
Result	R				S	T0 -				_	Status information after Command	Command	w		-	-	-	-	0			Command Codes
	R	_		•	_ s s	T2 _				=	execution		w		SPT. _ HL	T	-+-		'	HUT	→ND	Timer Settings
	R ·					н					Sector ID information after Command								STA			
	R										execution	Command	w	0	0	0	0	-	1 -	-	0 DS0	Command Codes
				SC	н и	IGH (OR E	QUA	AL_			Result	B		•	•	•				DS0	Status information
Command	w		MFM			1			-	1	Command Codes					,						about the FDD
	w		0	0	•	-	HD				Sector ID information							SEE	(
	w					н					prior Command	Command	w	0	0	-	-		1			Command Codes
1	w					R N		_			execution	1	w	0	0	-	-	-	HDS	D\$1	DS0	
	w				E	OT_						1	w					С _				ļ
	W										·	Execution										Head is positioned over proper Cylinder
Execution											Data compared		<u> </u>	L				<u>, </u>				on Diskette
	İ				•						between the FDD and the main-system							NVAL	.ID			,
Result	R				s	T1.					Status information after Command	Command	W			!	nvalio	l Cod	es _			Invalid Command Codes (NoOp FDC goes into Standby
,	R	i				¢					execution . Sector ID information	Result	R				_ s	то_				State) ST 0 = 80 (16)
	R R	1				R					after Command execution											1

Table 6. Command/Result Parameter Abbreviations

Symbol	Description	Symbol	Description
С	Cylinder Address. The currently selected cylinder address (0 to 76) on the disk.	EOT	End of Track. The final sector number of the current track.
D DS0,DS1	Data Pattern. The pattern to be written in each sector data field during formatting. Disk Drive Select.	GPL	Gap Length. The gap 3 size. (Gap 3 is the space between sectors excluding the VCO synchronization field as defined in section 3.)
D30,D31	DS1 DS0 0 0 Drive 0 0 1 Drive 1	н	Head Address. Selected head: 0 or 1 (disk side 0 or 1, respectively) as encoded in the sector ID field.
DTL	1 0 Drive 2 1 1 Drive 3 Special Sector Size. During the execution of disk read/write commands, this parameter is used to temporarily after the effective disk sec-	HLT ·	Head Load Time. Defines the time interval that the FDC waits after loading the head before initiating a read or write operation. Programmable from 2 to 254 milliseconds (in increments of 2 ms).
	tor size. By setting N to zero, DTL may be used to specify a sector size from 1 to 256 bytes in length. If the actual sector (on the diskette) is larger than DTL specifies, the remainder of the actual sector is not passed to	нит	Head Unload Time. Defines the time interval from the end of the execution phase (of a read or write command) until the head is unloaded. Programmable from 16 to 240 milliseconds (in increments of 16 ms).
	the system during read commands; during write commands, the remainder of the actual sector is written with all-zeroes bytes. DTL should be set to FF hexadecimal when N is not zero.	MFM	MFM/FM Mode Selector. Selects MFM double-density recording mode when high, FM single-density mode when low.

Table 6. Command/Result Parameter Abbreviations (continued)

Symbol	Description	Symbol	Description
MT	Multi-Track Selector. When set, this flag selects the multi-track operating mode. In this mode (used only with dual-sided disks), the FDC treats a complete cylinder (under both read/write head 0 and read/write head 1) as a single track. The FDC operates as if this expanded track started at the first sector under head 0 and ended at the last sector under head 1. With this flag set (high), a multi-sector read operation will automatically continue to the first sector under head 1 when the FDC finishes operating on the last sector under head 0.	SK SRT	Skip Flag. When this flag is set, sectors containing deleted data address marks will automatically be skipped during the execution of multi-sector Read Data or Scan commands. In the same manner, a sector containing a data address mark will automatically be skipped during the execution of a multi-sector Read Deleted Data command. Step Rate Interval. Defines the time interval between step pulses issued by the FDC (trackto-track access time). Programmable from 1 to 16 milliseconds (in increments of 1 ms).
N ND	Sector Size. The number of data bytes within a sector. (See Table 9.) Non-DMA Mode Flag. When set (high), this flag indicates that the FDC is to operate in the non-DMA mode. In this mode, the processor is interrupted for each data transfer. When low, the FDC interfaces to a DMA controller by means of the DRQ and DACK signals.	ST0 ST1 ST2 ST3	Status Register 0-3. Registers within the FDC that store status information after a command has been executed. This status information is available to the processor during the Result Phase after command execution. These registers may only be read after a command has been executed (in the exact order shown in Table 5 for each command). These registers should not be confused with the Main Status
R SC	Sector Address. Specifies the sector number to be read or written. In multi-sector transfers, this parameter specifies the sector number of the first sector to be read or written. Number of Sectors per Track. Specifies the	STP	Register. Scan Sector Increment: During Scan operations, this parameter is added to the current sector number in order to determine the next sector to be scanned.
50	number of sectors per track to be initialized by the Format Track command.		

Table 7. Status Register Definitions

Bit Number	Symbol	Description						
Status Register 0								
7,6	IC	Interrupt Code.						
	٠.	00 — Normal termination of command. The specified command was properly executed and completed without error.						
		01 — Abnormal termination of command. Command execution was started but could not be successfully completed.						
		10 — Invalid command. The requested command could not be executed.						
		11 — Abnormal termination. During command execution, the disk drive ready signal changed state.						
5	SE	Seek End. This flag is set (high) when the FDC has completed the Seek command and the read/write head is positioned over the correct cylinder.						
4	EC .	Equipment Check Error. This flag is set (high) if a fault signal is received from the disk drive or if the track 0 signal fails to become active after 77 step pulses (Recalibrate command).						
3	NR	Not Ready Error. This flag is set if a read or write command is issued and either the drive is not ready or the command specifies side 1 (head 1) of a single-sided disk.						
2	н	Head Address. The head address at the time of the interrupt.						
1,0	DS1,DS0	Drive Select. The number of the drive selected at the time of the interrupt.						

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Table 7. Status Register Definitions (continued)

Bit		Table 7. Status Register Definitions (continued)
Number	Symbol	Description
Status Regis	ster 1	
7	EN	End of Track Error. This flag is set if the FDC attempts to access a sector beyond the final sector of the track.
6		Not used. This bit is always low.
5	DE	Data Error. Set when the FDC detects a CRC error in either the ID field or the data field of a sector.
4	OR	Overrun Error. Set (during data transfers) if the FDC does not receive DMA or processor service within the specified time interval.
3	'	Not used. This bit is always low.
2	ND 	Sector Not Found Error. This flag is set by any of the following conditions. a) The FDC cannot locate the sector specified in the Read Data, Read Deleted Data, or Scan command. b) The FDC cannot locate the starting sector specified in the Read Track command.
1	NW	c) The FDC cannot read the ID field without error during a Read ID command. Write Protect Error. This flag is set if the FDC detects a write protect signal from the disk drive during the execution of a Write Data, Write Deleted Data, or Format Track command.
Ó	МА	Missing Address Mark Error. This flag is set by either of the following conditions: a) The FDC cannot detect the ID address mark on the specified track (after two occurrences of the physical index mark). b) The FDC cannot detect the data address mark or deleted data address mark on the specified track. (See also the MD bit of Status Register 2.)
Status Regis	iter 2	·
. 7		Not used. This bit is always low.
6	СМ	Control Mark. This flag is set when the FDC encounters one of the following conditions: a) A deleted data address mark during the execution of a Read Data or Scan command. b) A data address mark during the execution of a Read Deleted Data command.
. 5	DD	Data Error. Set (high) when the FDC detects a CRC error in a sector data field. This flag is not set when a CRC error is detected in the ID field.
4	wc	Cylinder Address Error. Set when the cylinder address from the disk sector ID field is different from the current cylinder address maintained within the FDC.
3	SH	Scan Hit. Set during the execution of the Scan command if the scan condition is satisfied.
2.	SN	Scan Not Satisfied. Set during execution of the Scan command if the FDC cannot locate a sector on the specified cylinder that satisfies the scan condition.
1	ВС	Bad Track Error. Set when the cylinder address from the disk sector ID field is FF hexadecimal and this cylinder address is different from the current cylinder address maintained within the FDC. This all "ones" cylinder number indicates a bad track (one containing hard errors) according to the IBM soft-sectored format specifications.
0	MD	Missing Data Address Mark Error. Set if the FDC cannot detect a data address mark or deleted data address mark on the specified track.

Table 7. Status Register Definitions (continued)

Bit Number	Symbol	Description
Status Regi	ister 3	
7	FT	Fault. This flag indicates the status of the fault signal from the selected disk drive.
6 ·	WP	Write Protected. This flag indicates the status of the write protect signal from the selected disk drive.
5	RDY	Ready. This flag indicates the status of the ready signal from the selected disk drive.
4	то	Track 0. This flag indicates the status of the track 0 signal from the selected disk drive.
3	TS	Two-Sided. This flag indicates the status of the two-sided signal from the selected disk drive.
2	Н	Head Address. This flag indicates the status of the side select signal for the currently selected disk drive.
1,0	DS1,DS0	Drive Select. Indicates the currently selected disk drive number.

In the non-DMA mode, transfer requests are indicated by activation of both the INT output signal and the RQM flag (bit 7) in the Main Status Register. INT can be used for interrupt-driven systems and RQM can be used for polled systems. The system processor must respond to the transfer request by reading data from (activating RD), or writing data to (activating WR), the FDC. This response removes the transfer request (INT and RQM are set inactive). After completing the last transfer, the 8272 activates the INT output to indicate the beginning of the result phase. In the non-DMA mode, the processor must activate the TC signal to the FDC (normally by means of an I/O port) after the transfer request for the last data byte has been received (by the processor) and before the appropriate data byte has been read from (or written to) the FDC.

In either mode of operation (DMA or non-DMA), the execution phase ends when a terminal count signal is sensed or when the last sector on a track (the EOT parameter—Table 5) has been read or written. In addition, if the disk drive is in a "not ready" state at the beginning of the execution phase, the "not ready" flag. (bit 3 in Status Register 0) is set (high) and the command is terminated.

If a fault signal is received from the disk drive at the end of a write operation (Write Data, Write Deleted Data, or Format), the FDC sets the "equipment check" flag (bit 4 in Status Register 0), and terminates the command after setting the interrupt code (bits 7 and 6 of Status Register 0) to "01" (bit 7 low, bit 6 high).

Multi-sector and Multi-track Transfers

During disk read/write transfers (Read Data, Write Data, Read Deleted Data, and Write Deleted Data), the FDC will continue to transfer data from sequential sectors until the TC input is sensed. In the DMA mode, the

TC input is normally connected to the TC/EOP (terminal count) output of the DMA controller. In the non-DMA mode, the processor directly controls the FDC TC input as previously described. Once the TC input is received, the FDC stops requesting data transfers (from the system processor or DMA controller). The FDC, however, continues to read data from, or write data to, the floppy disk until the end of the current disk sector. During a disk read operation, the data read from the disk (after reception of the TC input) is discarded, but the data CRC is checked for errors; during a disk write operation, the remainder of the sector is filled with all-zero bytes.

If the TC signal is not received before the last byte of the current sector has been transferred to/from the system, the FDC increments the sector number by one and initiates a read or write command for this new disk sector.

The FDC is also designed to operate in a multi-track mode for dual-sided disks. In the multi-track mode (specified by means of the MT flag in the command byte—Table 5) the FDC will automatically increment the head address (from 0 to 1) when the last sector (on the track under head 0) has been read or written. Reading or writing is then continued on the first sector (sector 1) of head 1.

Drive Status Polling

After the power-on reset, the 8272 automatically enters a drive status polling mode. If a change in drive status is detected (all drives are assumed to be "not ready" at power-on), an interrupt is generated. The 8272 continues this status polling between command executions (and between step pulses in the Seek command). In this manner, the 8272 automatically notifies the system processor when a floppy disk is inserted, removed, or changed by the operator.

Command Details

During the command phase, the Main Status Register must be polled by the CPU before each byte is written into the Data Register. The DIO (bit 6) and RQM (bit 7) flags in the Main Status Register must be low and high, respectively, before each byte of the command may be written into the 8272. The beginning of the execution phase for any of these commands will cause DIO to be set high and RQM to be set low.

The following paragraphs describe the fifteen FDC commands in detail.

Specify

The Specify command is used prior to performing any disk operations (including the formatting of a new disk) to define drive/FDC operating characteristics. The Specify command parameters set the values for three internal timers:

- Head Load Time (HLT) This seven-bit value defines the time interval that the FDC waits after loading the head before initiating a read or write operation. This timer is programmable from 2 to 254 milliseconds in increments of 2 ms.
- 2. Head Unload Time (HUT) This four-bit value defines the time from the end of the execution phase (of a read or write command) until the head is unloaded. This timer is programmable from 16 to 240 milliseconds in increments of 16 ms. If the processor issues another command before the head unloads, the head will remain loaded and the head load wait will be eliminated.
- Step Rate Time (SRT) This four-bit value defines the time interval between step pulses issued by the FDC (track-to-track access time). This timer is programmable from 1 to 16 milliseconds in increments of 1 ms.

The time intervals mentioned above are a direct function of the FDC clock (CLK on pin 19). Times indicated above are for an 8 MHz clock.

The Specify command also indicates the choice of DMA or non-DMA operation (by means of the ND bit). When this bit is high the non-DMA mode is selected; when ND is low, the DMA mode is selected.

Sense Drive Status

This command may be used by the processor whenever, it wishes to obtain the status of the disk drives. Status Register 3 (returned during the result phase) contains the drive status information as described in Table 7.

Sense Interrupt Status

An interrupt signal is generated by the FDC when one or more of the following events occurs:

- 1. The FDC enters the result phase for:
 - a. Read Data command
 - b. Read Track command
 - c. Read ID command
 - d. Read Deleted Data command
 - e. Write Data command
 - f. Format Track command
 - g. Write Deleted Data command
 - h. Scan commands
- The ready signal from one of the disk drives changes state.
- 3. A Seek or Recalibrate command completes operation.
- 4. The FDC requires a data transfer during the execution phase of a command in the non-DMA mode.

Interrupts caused by reasons (1) and (4) above occur during normal command operations and are easily discernible by the processor. However, interrupts caused by reasons (2) and (3) above are uniquely identified with the aid of the Sense Interrupt Status command. This command, when issued, resets the interrupt signal and by means of bits 5, 6, and 7 of Status Register 0 (returned during the result phase) identifies the cause of the interrupt (see Table 8).

Table 8. Interrupt Codes

		0	T .		
Seek End	Interrupt Code		Cause		
Bit 5	Bit 6	Bit 7			
0	1	1	Ready Line changed state, either polarity		
1	0	0	Normal Termination of Seek or Recalibrate Command		
1	1	0	Abnormal Termination of Seek or Recalibrate Command		

Neither the Seek nor the Recalibrate command has a result phase. Therefore, it is mandatory to use the Sense Interrupt Status Command after these commands to effectively terminate them and to provide verification of the disk head position.

When an interrupt is received by the processor, the FDC busy flag (bit 4) and the non-DMA flag (bit 5) may be used to distinguish the above interrupt causes:

bit	5	bit	4	

0	0	Asynchronous	event-(2)	or	(3) above
---	---	--------------	-----------	----	-----------

- 0 1 Result phase-(1) above
- 1 Data transfer required-(4) above

A single interrupt request to the processor may, in fact, be caused by more than one of the above events. The processor should continue to issue Sense Interrupt Status commands (and service the resulting conditions) until an invalid command code is received. In this manner, all "hidden" interrupts are serviced.

Seek

The Seek command causes the drive's read/write head to be positioned over the specified cylinder. The FDC determines the difference between the current cylinder address and the desired (specified) address, and issues the appropriate number of step pulses. If the desired cylinder address is larger than the current address, the direction signal (LCT/DIR, pin 38) is set high (step-in); the direction signal is set low (step-out) if the desired cylinder address is less than the current address. No head movement occurs (no step pulses are issued) if the desired cylinder is the same as the current cylinder.

The rate at which step pulses are issued is controlled by the step rate time (SRT) in the Specify command. After each step pulse is issued, the desired cylinder address is compared against the current cylinder address. When the cylinder addresses are equal, the "seek end" flag (bit 5 in Status Register 0) is set (high) and the command is terminated. If the disk drive becomes "not ready" during the seek operation, the "not ready" flag (in Status Register 0) is set (high) and the command is terminated.

During the command phase of the Seek operation the FDC is in the FDC busy state, but during the execution phase it is in the non-busy state. While the FDC is in the non-busy state, another Seek command may be issued. In this manner parallel seek operations may be in operation on up to four floppy disk drives at once. The Main Status Register contains a flag for each drive (Table 4) that indicates whether the associated drive is currently operating in the seek mode. When a drive has completed a seek operation, the FDC generates an interrupt. In response to this interrupt, the system software must issue a Sense Interrupt Status command. During the result phase of this command, Status Register 0 (containing the drive number in bits 0 and 1) is read by the processor.

Recalibrate

This command causes the read/write head of the disk drive to retract to the track 0 position. The FDC clears the contents of its internal cylinder counter, and checks the status of the track 0 signal from the disk drive. As long as the track 0 signal is low, the direction signal remains high and step pulses are issued. When the track 0 signal goes high, the seek end flag (in Status Register 0) is set (high) and the command is terminated. If the track 0 signal is still low after 77 step pulses have been issued, the seek end and equipment check flags (in Status Register 0) are both set and the Recalibrate command is terminated.

Recalibrate commands for multiple drives can be overlapped in the same manner that Seek commands are overlapped.

Format Track

The Format Track command formats or "initializes" a track on a floppy disk by writing the ID field, gaps, and address marks for each sector. Before issuing the Format command, the Seek command must be used to position the read/write head over the correct cylinder. In addition, a table of ID field values (cylinder, head, and sector addresses and sector length code) must be prepared before the command is executed. During command execution, the FDC accesses the table and, using the values supplied, writes each sector on the track. The ID field address mark originates from the FDC and is written automatically as the first byte of each sector's ID field. The cylinder, head, and sector addresses are taken, in order, from the table. The ID field CRC character (derived from the data written in the first five bytes) is written as the last two bytes of the ID field. Gaps are written automatically by the FDC, with the length of the variable gap determined by one of the Format command parameters.

The data field address mark is generated by the FDC and is written automatically as the first byte of the data field. The data pattern specified in the command phase is written into each data byte of each sector. A CRC character is derived from the data address mark and the data written in the sector's data field. The two CRC bytes are appended to the last data byte.

The formatting of a track begins at the physical index mark. As previously mentioned, the order of sector assignment is taken directly from the formatting table. Four entries are required for each sector: a cylinder address, a head address, a sector address, and a sector length code. The cylinder address in the ID field should be equal to the cylinder address of the track currently being formatted.

The sector addresses must be unique (no two equal). The order of the sector entries in the table is the sequence in which sector numbers appear on the track when it is formatted. The number of entry sets (cylinder, head, and sector address and sector length code) must equal the number of sectors allocated to the track (specified in the command phase).

Since the sector address is supplied, in order, for each sector, tracks can be formatted sequentially (the first sector following the index mark is assigned sector address 1, the adjacent sector is assigned sector address 2, and so on) or sector numbers can be interleaved (see section 3) on a track.

Table 9 lists recommended gap sizes and sectors/track for various sector sizes.

Read Data

Nine (9) bytes are required to complete the command phase specification for the Read Data command. During the execution phase, the FDC loads the head (if it is in the unloaded state), waits the specified head load time (defined in the Specify command), and begins reading ID address marks and ID fields. When the requested sector address compares with the sector address read from the disk, the FDC outputs data (from the data field) byte-by-byte to the system. The Read Data command automatically operates in the multi-sector mode described earlier. In addition, multi-track operation may be specified by means of the MT command flag (Table 5). The amount of data that can be transferred with a single command to the FDC depends on the multi-track flag, the recording density flag, and the number of bytes per sector.

During the execution of read and write commands, the special sector size parameter (DTL) is used to temporarily alter the effective disk sector size. By setting the sector size code (N) to zero, DTL may be used to specify a sector size from 1 to 256 bytes in length. If the actual sector (on the disk) is larger than DTL specifies, only the number of bytes specified by the DTL parameter are

passed to the system; the remainder of the actual disk sector is not transferred (although the data is checked for CRC errors). Multi-sector read operations are performed in the same manner as they are when the sector size code is non-zero. (The N and DTL parameters are always present in the command sequence. DTL should be set to FF hexadecimal when N is not zero.)

If the FDC detects the physical index mark twice without finding the requested sector, the FDC sets the "sector not found error" flag (bit 2 in Status Register 1) and terminates the Read Data command. The interrupt code (bits 7 and 6 of Status Register 0) is set to "01." Note that the FDC searches for each sector in a multisector operation. Therefore, a "sector not found" error may occur after successful transfer of one or more preceding sectors. This error could occur if a particular sector number was not included when the track was first formatted or if a hard error on the disk has invalidated a sector ID field.

After reading the ID field and data field in each sector, the FDC checks the CRC bytes. If a read error is detected (incorrect CRC in the ID field), the FDC sets the "data error" flag in Status Register 1; if a CRC error occurs in the data field, the FDC sets the "data error" flag in Status Register 2. In either error condition, the FDC terminates the Read Data command. The interrupt code (bits 7 and 6 in Status Register 0) is set to "01."

If the FDC reads a deleted data address mark from the disk, and the skip flag (specified during the command phase) is not set, the FDC sets the "control mark" flag (bit 6 in Status Register 2) and terminates the Read Data command (after reading all the data in the sector). If the skip flag is set, the FDC skips the sector with the deleted data address mark and reads the next sector. Thus, the skip flag may be used to cause the FDC to ignore deleted data sectors during a multi-sector read operation.

During disk data transfers between the FDC and the system, the FDC must be serviced by the system (processor or DMA controller) every $27 \mu s$ in the FM mode, and every $13 \mu s$ in the MFM mode. If the FDC is not

Table 9. Sector Size Relationships

Format	Sector Size	N Sector Size Code	SC Sectors/ Track	GPL ¹ Gap 3 Length	GPL ² Gap 3 Length	Remarks
FM Mode	128 bytes/Sector 256 512	00 01 02	1A ₍₁₆₎ 0F ₍₁₆₎ 08	07 ₍₁₆₎ 0E ₍₁₆₎ 1B ₍₁₆₎	1B ₍₁₆₎ 2A ₍₁₆₎ 3A ₍₁₆₎	IBM Diskette 1 IBM Diskette 2
MFM Mode	256 512 1024	01 02 03	1A ₍₁₆₎ 0F ₍₁₆₎ 08	0E ₍₁₆₎ 1B ₍₁₆₎ 35 ₍₁₆₎	36 ₍₁₆₎ 54 ₍₁₆₎ 74 ₍₁₆₎	IBM Diskette 2D IBM Diskette 2D

Notes: 1. Suggested values of GPL in Read or Write commands to avoid splice point between data field and ID field of contiguous sectors.

2. Suggested values of GPL in Format command.

serviced within this interval, the "overrun error" flag (bit 4 in Status Register 1) is set and the Read Data command is terminated.

If the processor terminates a read (or write) operation in the FDC, the ID information in the result phase is dependent upon the state of the multi-track flag and end of track byte. Table 11 shows the values for C, H, R, and N, when the processor terminates the command.

Write Data

Nine (9) bytes are required to complete the command phase specification for the Write Data command. During the execution phase the FDC loads the head (if it is in the unloaded state), waits the specified head load time (defined by the Specify command), and begins reading sector ID fields. When the requested sector address compares with the sector address read from the disk, the FDC reads data from the processor one byte at a time via the data bus and outputs the data to the data field of that sector. The CRC is computed on this data and two CRC bytes are written at the end of the data field.

The FDC reads the ID field of each sector and checks the CRC bytes. If the FDC detects a read error (incorrect CRC) in one of the ID fields, it sets the "data error" flag (bit 5 in Status Register 1) and terminates the Write Data command. The interrupt code (bits 7 and 6 in Status Register 0) is set to "01."

The Write Data command operates in much the same manner as the Read Data command. The following items are the same; refer to the Read Data command for details:

- Multi-sector and Multi-track operation
- Data transfer capacity
- "End of track error" flag
- · "Sector not found error" flag
- · "Data error" flag
- · Head unload time interval
- ID information when the processor terminates the command (see Table 11)
- Definition of DTL when N=0 and when $N\neq 0$

During the Write Data execution phase, data transfers between the processor and FDC must occur every 31 μ s in the FM mode, and every 15 μ s in the MFM mode. If the time interval between data transfers is longer than this, the FDC sets the "overrun error" flag (bit 4 in Status Register 1) and terminates the Write Data command.

Read Deleted Data

This command operates in almost the same manner as the Read Data command operates. The only difference involves the treatment of the data address mark and the skip flag. When the FDC detects a data address mark at the beginning of a data field (and the skip flag is not set), the FDC reads all the data in the sector, sets the "control mark" flag (bit 6 in Status Register 2), and terminates the command. If the skip flag is set, the FDC skips the sector with the data address mark and continues reading at the next sector. Thus, the skip flag may be used to cause the FDC to read only deleted data sectors during a multi-sector read operation.

Write Deleted Data

This command operates in the same manner as the Write Data command operates except that a deleted data address mark is written at the beginning of the data field instead of the normal data address mark. This command is used to mark a bad sector (containing a hard error) on the floppy disk.

Read Track

The Read Track command is similar to the Read Data command except that the entire data field is read continuously from each of the sectors of a track. Immediately after encountering the physical index mark, the FDC starts reading all data fields on the track as continuous blocks of data. If the FDC finds an error in the ID field or data field CRC check bytes, it continues to read data from the track. The FDC compares the ID information read from each sector with the values specified during the command phase. If the specified ID field information is not found on the track, the "sector not found error" flag (in Status Register 1) is set. Multitrack and skip operations are not allowed with this command.

This command terminates when the last sector on the track has been read. (The number of sectors on the track is specified by the end of track parameter byte during the command phase.) If the FDC does not find an ID address mark on the disk after it encounters the physical index mark for the second time, it sets the "missing address mark error" flag (bit 0 in Status Register 1) and terminates the command. The interrupt code (bits 7 and 6 of Status Register 0) is set to "01."

Read ID

The Read ID command transfers (reads) the first correct ID field from the current disk track (following the physical index mark) to the processor. If no correct ID address mark is found on the track, the "missing address mark error" flag is set (bit 0 in Status Register 1). If no data mark is found on the track, the "sector not found error" flag is also set (bit 2 in Status Register 1). Either error condition causes the command to be terminated.

Scan Commands

The Scan commands allow the data being read from the disk to be compared against data supplied by the system (by the processor in non-DMA mode, and by the DMA controller in DMA mode). The FDC compares the data on a byte-by-byte basis, and searches for a sector of data that meets the conditions of "disk data equal to system data", "disk data less than or equal to system data", or "disk data greater than or equal to system data". Simple binary (ones complement) arithmetic is used for comparison (FF = largest number, 00 = smallest number). If, after a complete sector of data is compared, the conditions are not met, the sector number is incremented by the scan sector increment (specified in the command phase), and the scan operation is continued. The scan operation continues until one of the following conditions occurs; the conditions for scan are met (equal, low, or high), the last sector on the track is reached, or the terminal count signal is received.

If the conditions for scan are met, the FDC sets the "scan hit" flag (bit 3 in Status Register 2) and terminates the Scan command. If the conditions for scan

are not met between the starting sector and the last sector on the track (specified in the command phase), the FDC sets the "scan not satisfied" flag (bit 2 in Status Register 2) and terminates the Scan command. The receipt of a terminal count signal from the processor or DMA controller during the scan operation will cause the FDC to complete the comparison of the particular byte which is in process, and to terminate the command. Table 10 shows the status of the "scan hit" and "scan

Table 10. Scan Status Codes

Command	Status R	egister 2	Comments		
00	Bit 2 = SN Bit 3 = SH				
Scan Equal	0 1	. 1 0	D _{FDD} =D _{Processor} D _{FDD} ≠ D _{Processor}		
Scan Low or Equal	0 0 1	1 0 0	D _{FDD} = D _{Processor} D _{FDD} < D _{Processor} D _{FDD} ≰ D _{Processor}		
Scan High or Equal	0 0 1	1 0 0	$D_{FDD} = D_{Processor}$ $D_{FDD} > D_{Processor}$ $D_{FDD} \not \ge D_{Processor}$		

Table 11. ID Information When Processor Terminates Command

мт ЕОТ	EOT	Final Sector Transferred to		ID Information at Result Phase				
	Processor	С	н	R	N			
	1A 0F 08	Sector 1 to 25 at Side 0 Sector 1 to 14 at Side 0 Sector 1 to 7 at Side 0	NC	NC .	R+1	NC		
0	1A 0F 08	Sector 26 at Side 0 Sector 15 at Side 0 Sector 8 at Side 0	C+1	NC	R=01	NC		
	1A 0F 08	Sector 1 to 25 at Side 1 Sector 1 to 14 at Side 1 Sector 1 to 7 at Side 1	NC #	NC	R+1	NC		
	1A 0F 08	Sector 26 at Side 1 Sector 15 at Side 1 Sector 8 at Side 1	C+1	NC	R = 01	NC		
	1A 0F 08	Sector 1 to 25 at Side 0 Sector 1 to 14 at Side 0 Sector 1 to 7 at Side 0	NC	NC	R+1	NC		
1	1A 0F 08	Sector 26 at Side 0 Sector 15 at Side 0 Sector 8 at Side 0	NC	LSB	R = 01	NC		
	1A 0F 08	Sector 1 to 25 at Side 1 Sector 1 to 14 at Side 1 Sector 1 to 7 at Side 1	NC	NC	R+1	NC		
	1A 0F 08	Sector 26 at Side 1 Sector 15 at Side 1 Sector 8 at Side 1	C+1	LSB	R=01	NC		

Notes: 1. NC (No Change): The same value as the one at the beginning of command execution.

2. LSB (Least Significant Bit): The least significant bit of H is complemented.

not satisfied" flags under various scan termination conditions.

If the FDC encounters a deleted data address mark in one of the sectors and the skip flag is low, it regards the sector as the last sector on the cylinder, sets the "control mark" flag (bit 6 in Status Register 2) and terminates the command. If the skip flag is high, the FDC skips the sector with the deleted address mark, and reads the next sector. In this case, the FDC also sets the "control mark" flag (bit 6 in Status Register 2) in order to show that a deleted sector had been encountered.

NOTE: During scan command execution, the last sector on the track must be read for the command to terminate properly. For example, if the scan sector increment is set to 2, the end of track parameter is set to 26, and the scan begins at sector 21, sectors 21, 23, and 25 will be scanned. The next sector, 27 will not be found on the track and an abnormal command termination will occur. The command would be completed in a normal manner if either a) the scan had started at sector 20 or b) the end of track parameter had been set to 25.

During the Scan command, data is supplied by the processor or DMA controller for comparison against the data read from the disk. In order to avoid having the "overrun error" flag set (bit 4 in Status Register 1), it is necessary to have the data available in less than 27 μ s (FM Mode) or 13 μ s (MFM Mode). If an overrun error occurs, the FDC terminates the command.

Invalid Commands

If an invalid (undefined) command is sent to the FDC, the FDC will terminate the command. No interrupt is generated by the 8272 during this condition. Bit 6 and bit 7 (DIO and RQM) in the Main Status Register are both set indicating to the processor that the 8272 is in the result phase and the contents of Status Register 0 must be read. When the processor reads Status Register 0 it will find an 80H code indicating that an invalid command was received.

A Sense Interrupt Status command must be sent after a Seek or Recalibrate interrupt; otherwise the FDC will consider the next command to be an invalid command. Also, when the last "hidden" interrupt has been serviced, further Sense Interrupt Status commands will result in invalid command codes.

In some applications the user may wish to use this command as a No-Op command to place the FDC in a stand-by or no operation state.

5. THE DATA SEPARATOR

As briefly discussed in section 2, LSI disk controllers such as the 8272 require external circuitry to generate a data window signal. This signal is used within the FDC to isolate the data bits contained within the READ DATA input signal from the disk drive. (The disk READ DATA signal is a composite signal constructed from both clock and data information.) After isolating the data bits from this input signal, the FDC assembles the data bits into 8-bit bytes for transfer to the system processor or memory.

Single Density

In single-density (FM) recording (Figure 3), the bit cell is 4 microseconds wide. Each bit cell contains a clock bit at the leading edge of the cell. The data bit (if present) is always located at the center of the cell. The job of data separation is relatively straightforward for single-density; simply generate a data window $2 \mu s$ wide starting $1 \mu s$ after each clock bit. Since every cell has a clock bit, a fixed window reference is available for every data bit and because the window is $2 \mu s$ wide, a slightly shifted data bit will still remain within the data window.

A single-density data separator with these specifications may be easily generated using a digital or analog oneshot triggered by the clock bit.

Double-Density

Double-density (MFM) bit cells are reduced to 2 μ s (in order to double the disk data storage capacity). Clock bits are inserted into the data stream only if data bits are not present in both the current and preceding bit cells (Figure 3). The data bit (if present) still occurs at the center of the bit cell and the clock bit (if present) still occurs at the leading edge of the bit cell.

MFM data separation has two problems. First, only some bit cells contain a clock bit. In this manner, MFM encoding loses the fixed bit cell reference pulse present in FM encoding. Second, the bit cell for MFM is one-half the size of the bit cell for FM. This shorter bit cell means that MFM cannot tolerate as large a playback data-shift (as FM can tolerate) without errors.

Since most playback data-shift is predictable, the FDC can precompensate the write data stream so that data/clock pulses will be correctly positioned for subsequent playback. This function is completely controlled by the FDC and is only required for MFM recording. During write operations, the FDC specifies an early, normal, or late bit positioning. This timing information is specified with respect to the FDC write clock. Early and late timing is typically 125 ns to 250 ns before or after the write clock transition (depending on disk drive requirements).

The data separator circuitry for double-density recording must continuously analyze the total READ DATA stream, synchronizing its operation (window generation) with the actual clock/data bits of the data stream. The data separation circuit must track the disk input data frequency very closely—unpredictable bit shifts leave less than 50 ns margin to the window edges.

Phase-Locked Loop

Only an analog phase-locked loop (PLL) can provide the reliability required for a double-density data separation circuit. (A phase-locked loop is an electronic circuit that constantly analyzes the frequency of an input signal and locks another oscillator to that frequency.) Using analog PLL techniques, a data separator can be designed with ±1 ns resolution (this would require a 100 MHz clock in a digital phase-locked loop). The analog PLL determines the clock and data bit positions by sampling each bit in the serial data stream. The phase relationship between a data bit and the PLL generated data window is constantly fed back to adjust the position of the data window, enabling the PLL to track input data frequency changes, and thereby reliably read previously recorded data from a floppy disk.

PLL Design

A block diagram of the phase-locked loop described in this application note is shown in Figure 7. Basically, the phase-locked loop operates by comparing the frequency of the input data (from the disk drive) against the frequency of a local oscillator. The difference of these frequencies is used to increase or decrease the frequency of the local oscillator in order to bring its frequency closer to that of the input. The PLL synchronizes the local oscillator to the frequency of the input during the all "zeroes" synchronization field on the floppy disk (immediately preceding both the ID field and the data field).

The PLL consists of nine ICs and is located on page 3 of the schematics in the Appendix. The 8272 VCO output essentially turns the PLL circuitry on and off. When the PLL is off, it "idles" at its center frequency. The VCO output turns the PLL on only when valid data is being received from the disk drive. The VCO turns the PLL on after the read/write head has been loaded and the head load time has elapsed. The PLL is turned off in the gap between the ID field and the data field and in the gap after the data field (before the next sector ID field). The GPL parameter in the FDC read and write commands specifies the elapsed time (number of data bytes) that the PLL is turned off in order to blank out discontinuities that appear in the gaps when the write current is turned on and off. The PLL operates with either MFM or FM input data. The MFM output from the FDC controls the PLL operation frequency.

The PLL consists of six functional blocks as follows:

- Pulse Shaping A 96LS02 senses a READ DATA pulse and provides a clean output signal to the FDC and to the PLL Phase Comparator and Frequency Discriminator circuitry.
- 2. Phase Comparator The phase difference between the PLL oscillator and the READ DATA input is compared. Pump up (PU) and pump down (PD) error signals are derived from this phase difference and output to the filter. If there is no phase difference between the PLL oscillator and the READ DATA input, the PU and PD pulse widths are equal. If the READ DATA pulse occurs early, the PU duration is shorter than the PD duration. If the data pulse occurs late, the PU duration is longer than the PD duration.
- Filter This analog circuit filters the PU and PD pulses into an error voltage. This error voltage is buffered by an LM358 operational amplifier.

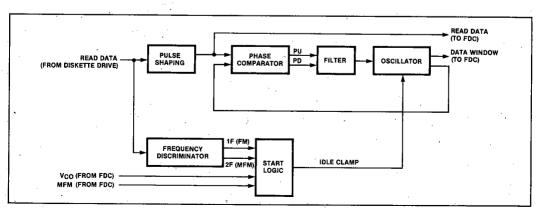


Figure 7. Phase-Locked Loop Data Separator

- PLL Oscillator This oscillator is composed of a 74LS393, 74LS74, and 96LS02. The oscillator frequency is controlled by the error voltage output by the filter. This oscillator also generates the data window signal to the FDC.
- Frequency Discriminator This logic tracks the READ DATA input from the disk drive and discriminates between the synchronization gap for FM recording (250 KHz) and the gap for MFM recording (500 KHz). Synchronization gaps immediately precede address marks.
- 6. Start Logic The function of this logic is to clamp the PLL oscillator to its center frequency (2 MHz) until the FDC VCO signal is enabled and a valid data pattern is sensed by the frequency discriminator. The start logic (consisting of a 74LS393 and 74LS74) ensures that the PLL oscillator is started with zero phase error.

PLL Adjustments

The PLL must be initially adjusted to operate at its center frequency with the VCO output off and the adjustment jumper removed. The 5K trimpot should be adjusted until the frequency at the test point (Q output of the 96LS02) is 2 MHz. The jumper should then be replaced for normal operation.

PLL Design Details

The following paragraphs describe the operational and design details of the phase-locked loop data separator il-

lustrated in the appendix. Note that the analog section is operated from a separately filtered +5V supply.

Initialization

As long as the 8272 maintains a low VCO signal, the data separator logic is "turned off". In this state, the PLL oscillator (96LS02) is not oscillating and therefore the 2XBR signal is constantly low. In addition, the pump up (PU) and pump down (PD) signals are inactive (PU low and PD high), the CNT8 signal is inactive (low), and the filter input voltage is held at 2.5 volts by two 1Mohm resistors between ground and +5 volts.

Floppy Disk Data

The data separator frequency discriminator, the input pulse shaping circuitry, and the start logic are always enabled and respond to rising edges of the READ DATA signal. The rising edge of every data bit from the disk drive triggers two pulse shaping one-shots. The first pulse shaper generates a stable and well-defined 200 ns read data pulse for input to the 8272 and other portions of the data separator logic. The second one-shot generates a 2.5 μ s data pulse that is used for input data frequency discrimination.

The frequency discriminator operates as illustrated in Figure 8. The 2F output signal is active (high) during reception of valid MFM (double-density) sync fields on the disk while the 1F signal is active (high) during reception of valid FM (single-density) sync fields. A multiplexer (controlled by the 8272 MFM signal) selects the appropriate 1F or 2F signal depending on the programmed mode.

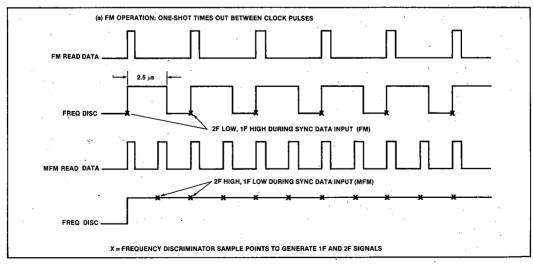


Figure 8. Input Data Frequency Discrimination

Startup

The data separator is designed to require reception of eight valid sync bits (one sync byte) before enabling the PLL oscillator and attempting to synchronize with the input data stream (see Figure 9). This delay ensures that the PLL will not erroneously synchronize outside a valid sync field in the data stream if the VCO signal is enabled slightly early. The sync bit counter is asynchronously reset by the CNTEN signal when valid sync data is not being received by the drive.

Once the VCO signal is active and eight sync bits have been counted, the CNT8 signal is enabled. This signal turns on the PLL oscillator. Note that this oscillator starts synchronously with the rising edge of the disk input data (because CNT8 is synchronous with the data rising edge) and the oscillator also starts at its center frequency of 2 MHz (because the LM348 filter input is held at its center voltage of approximately 2.5 volts). This frequency is divided by two and four to generate the 2XBR signal (1 MHz for MFM and 500 KHz for FM).

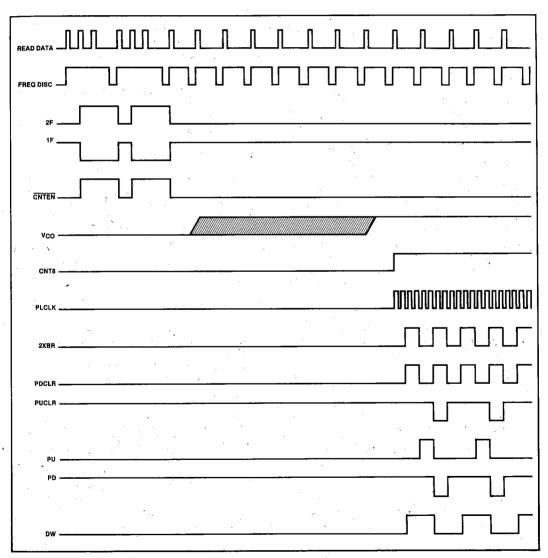


Figure 9. Typical Data Separator Startup Timing Diagram

PLL Synchronization

At this point, the PLL is enabled and begins to synchronize with the input data stream. This synchronization is accomplished very simply in the following manner. The pump up (PU) signal is enabled on the rising edge of the READ DATA from the disk drive. (When the PLL is synchronized with the data stream, this point will occur at the same time as the falling edge of the 2XBR signal as shown in Figure 9). The PU signal is turned off and the PD signal is activated on the next rising edge of the 2XBR clock. With this scheme, the difference between PU active time and the PD active time is equal to the difference between the input bit rate and the PLL clock rate. Thus, if PU is turned on longer than PD is on, the input bit rate is faster than the PLL clock.

As long as PU and PD are both inactive, no charge is transferred to or from the LM358 input holding capacitor, and the PLL output frequency is maintained (the LM358 operational amplifier has a very high input impedance). Whenever PU is turned on, current flows from the +5 volt supply through a 20K resistor into the holding capacitor. When the PD signal is turned on, current flows from the holding capacitor to ground through a 20K resistor. In this manner, both the pump up and pump down charging rates are balanced.

The change in capacitor charge (and therefore voltage) after a complete PU/PD cycle is proportional to the difference between the PU and PD pulse widths and is also proportional to the frequency difference between the incoming data stream and the PLL oscillator. As the capacitor voltage is raised (PU active longer than PD), the PLL oscillator time constant (RC of the 96LS02) is modified by the filter output (LM358) to raise the oscillator frequency. As the capacitor voltage is lowered (PD active longer than PD), the oscillator frequency is lowered. If both frequencies are equal, the voltage on the holding capacitor does not change, and the PLL oscillator frequency remains constant.

6. AN INTELLIGENT DISKETTE DATA BASE SYSTEM

The system described in this application note is designed to function as an intelligent data base controller. The schematics for this data base unit are presented in Appendix A; a block diagram of the unit is illustrated in Figure 10. As designed, the unit can access over four million bytes of mass storage on four floppy disk drives (using a single 8272 FDC); the system can easily be expanded to four FDC devices (and 16 megabytes of online disk storage). Three serial data links are also included. These data links may be used by CRT terminals or other microprocessor systems to access the data base.

Processor and Memory

A high-performance 8088 eight-bit microprocessor (operating at 5 MHz with no wait states) controls system operation. The 8088 was selected because of its memory addressing capabilities and its sophisticated string handling instructions. These features improve the speed of data base search operations. In addition, these capabilities allow the system to be easily upgraded with additional memory, disk drives, and if required, a bubble memory or winchester disk unit.

The schematics for the basic design provide 8K bytes of 2732A high-speed EPROM program storage and 8K bytes of disk directory and file buffer RAM. This memory can easily be expanded to 1 megabyte for performance upgrades.

An 8259A Programmable Interrupt Controller (PIC) is also included in the design to field interrupts from both the serial port and the FDC. This interrupt controller provides a large degree of programming flexibility for the implementation of data base functions in an asynchronous, demand driven environment. The PIC allows the system to accumulate asynchronous data base requests from all serial I/O ports while previously specified data base operations are currently in progress. This feature is made possible by the ability of the 8251A RXRDY signal to cause a processor interrupt. After receiving this interrupt, the processor can temporarily halt work on existing requests and enter the incoming information into a data base request buffer. Once the information has been entered into the buffer, the system can resume its previous processing.

In addition, the PIC permits some portions of data base requests to be processed in parallel. For example, once a disk record has been loaded into a memory buffer, a memory search can proceed in parallel with the loading of the next record. After the FDC completes the record transfer, the memory search will be interrupted and the processor can begin another disk transfer before resuming the memory search.

The bus structure of the system is split into three functional buffered units. A 20-bit address from the processor is latched by three-state transparent 74LS373 devices. When the processor is in control of the address and data busses, these devices are output enabled to the system buffered address bus. All I/O devices are placed directly on the local data bus. Finally, the memory data bus is isolated from the local data bus by an 8286 octal transceiver. The direction of this transceiver is determined by the Memory Read signal, while its output enable is activated by a Memory Read or Memory Write command.

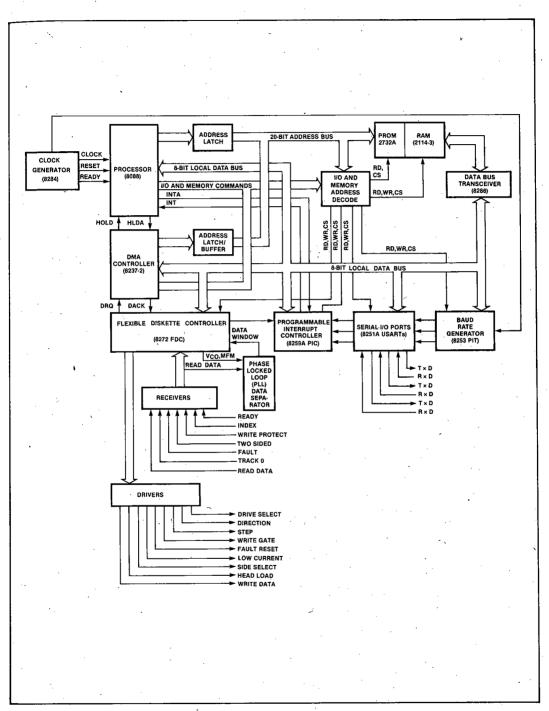


Figure 10. Intelligent Data Base Block Diagram

Serial I/O

The three RS-232-C compatible serial I/O ports operate at software-programmable baud rates to 19.2K. Each I/O port is controlled by an 8251A USART (Universal Synchronous/Asynchronous Receiver/Transmitter). Each USART is individually programmable for operation in many synchronous and asynchronous serial data transmission formats (including IBM Bi-sync). In operation, USART error detection circuits can check for parity, data overrun, and framing errors. An 8253 Programmable Interval Timer is employed to generate the baud rates for the serial I/O ports.

The Transmitter Ready and Receiver Ready output signals of the 8251As are routed to the interrupt inputs of the 8259A interrupt controller. These signals interrupt processor execution when a data byte is received by a USART and also when the USART is ready to accept another data byte for transmission.

DMA

The 8272 FDC interfaces to system memory by means of an 8237-2 high-speed DMA controller. Transfers between the disk controller and memory also operate with no wait states when 2114-3 (150 ns) or faster static RAM is used. In operation, the 8272 presents a DMA request to the 8237 for every byte of data to be transferred. This request causes the 8273 to present a HOLD request to the 8088. As soon as the 8088 is able to relinquish data/address bus control, the processor signals a HOLD acknowledge to the 8237. The 8237 then assumes control over the data and address busses. After latching the address for the DMA transfer, the 8237 generates simultaneous I/O Read and Memory Write commands (for a disk read) or simultaneous I/O Write and Memory Read commands (for a disk write). At the same time, the 8272 is selected as the I/O device by means of the DMA acknowledge signal from the 8237. After this single byte has been transferred between the FDC and memory, the DMA controller releases the data/address busses to the 8088 by deactivating the HOLD request. In a short period of time (13 μ s for double-density and 27 μs for single-density) the FDC requests a subsequent data transfer. This transfer occurs in exactly the same manner as the previous transfer. After all data transfers have been completed (specified by the word count programmed into the 8237 before the FDC operation was initiated), the 8237 signals a terminal count (EOP pin). This terminal count signal informs the 8272 that the data transfer is complete. Upon reception of this terminal count signal, the 8272 halts DMA requests and initiates an "operation complete" interrupt.

Since the system is designed for 20-bit addressing, a four-bit DMA-address latch is included as a processor

addressable I/O port. The processor writes the upper four DMA address bits before a data transfer. When the DMA controller assumes bus control, the contents of this latch are output enabled on the upper four bits of the address bus. The only restriction in the use of this address latch is that a single disk read or write transfer cannot cross a 64K memory boundary.

Disk Drive Interface

The 8272 FDC may be interfaced to a maximum of four eight-inch floppy disk drives. Both single- and double-density drives are accommodated using the data separation circuit described in section 5. In addition, single- or dual-sided disk drives may be used. The 8272 is driven by an 8 MHz crystal controller clock produced by an 8224 clock generator.

Drive select signals are decoded by means of a 74LS139 from the DS0, DS1 outputs of the FDC. The fault reset, step, low current, and direction outputs to the disk drives are generated from the FR/STEP, LCT/DIR, and RW/SEEK FDC output signals by means of a 74LS240. The other half of the 74LS240 functions as an input multiplexer for the disk write protect, two-sided, fault, and track zero status signals. These signals are multiplexed into the WP/TS and FLT/TRK0 inputs to the 8272.

The 8272 write clock (WR CLK) is generated by a ring counter/multiplexer combination. The write clock frequency is 1 MHz for MFM recording and 500 KHz for FM recording (selected by the MFM output of the 8272). The pulse width is a constant 250 ns. The write clock is constantly generated and input to the FDC (during both read and write operations). The FDC write enable output (WE) is transmitted directly to the write gate disk drive input.

Write data to the disk drive is preshifted (according to the PSO, PS1 FDC outputs) by the combination of a 74LS175 four-bit latch and a 74LS153 multiplexer. The amount of preshift is completely controlled within the 8272 FDC. Three cases are possible: the data may be written one clock cycle early, one clock cycle late, or with no preshift. The data preshift circuit is activated by the FDC only in the double-density mode. The preshift is required to cancel predictable playback data shifts when recorded data is later read from the floppy disk.

A single 50-conductor flat cable connects the board to the floppy disk drives. FDC outputs are driven by 7438 open collector high-current line-drivers. These drivers are resistively terminated on the last disk drive by means of a 150 ohm resistor to +5V. The line receivers are 7414 Schmitt triggered inverters with 150 ohm pull-up resistors on board.

7. SPECIAL CONSIDERATIONS

This section contains a quick review of key features and issues, most of which have been mentioned in other sections of this application note. Before designing with the 8272 FDC, it is advisable that the information in this section be completely understood.

1. Multi-Sector Transfers

The 8272 always operates in a multi-sector transfer mode. The 8272 continues to transfer data until the TC input is activated. In a DMA configuration, the TC input of the 8272 must always be connected to the EOP/TC output of the DMA controller. When multiple DMA channels are used on a single DMA controller, EOP must be gated with the select signal for the proper FDC. If the TC signal is not gated, a terminal count on another channel will abort FDC operation.

In a processor driven configuration with no DMA controller, the system must count the transfers and supply a TC signal to the FDC. In a DMA environment, ORing a programmable TC with the TC from the DMA controller is a convenient means of ensuring that the processor may always gain control of the FDC (even if the diskette system hangs up in an abnormal manner).

2. Processor Command/Result Phase Interface

In the command phase, the processor must write the exact number of parameters in the exact order shown in Table 5. During the result phase, the processor must read the complete result status. For example, the Format Track command requires six command bytes and presents seven result bytes. The 8272 will not accept a new command until all result bytes are read. Note that the number of command and result bytes varies from command-to-command. Command and result phases cannot be shortened.

During both the command and result phases, the Main Status Register must be read by the processor before each byte of information is read from, or written to, the FDC Data Register. Before each command byte is written, DIO (bit 6) must be low (indicating a data transfer from the processor) and RQM (bit 7) must be high (indicating that the FDC is ready for data). During the result phase, DIO must be high (indicating a data transfer to the processor) and RQM must also be high (indicating that data is ready for the processor).

NOTE: After the 8272 receives a command byte, the RQM flag may remain set for 12 microseconds (with an 8 MHz clock). Software should not attempt to read the Main Status Register before this time interval has elapsed; otherwise, the software will erroneously assume that the FDC is ready to accept the next byte.

3. Sector Sizes

The 8272 does not support 128 byte sectors in the MFM (double-density) mode.

4. Write Clock

The FDC Write Clock input (WR CLK) must be present at all times.

5. Reset

The FDC Reset input (RST) must be held active during power-on reset while the RD and WR inputs are active. If the reset input becomes inactive while RD and WR are still active, the 8272 enters the test mode. Once activated, the test mode can only be deactivated by a power-down condition.

6. Drive Status

The 8272 constantly polls (starting after the power-on reset) all drives for changes in the drive ready status. At power-on, the FDC assumes that all drives are not ready. If a drive application requires that the ready line be strapped active, the FDC will generate an interrupt immediately after power is applied.

7. Gap Length

Only the gap 3 size is software programmable. All other gap sizes are fixed. In addition, different gap 3 sizes must be specified in format, read, write, and scan commands. Refer to Section 3 and Table 9 for gap size recommendations.

8. Seek Command

The drive busy flag in the Main Status Register remains set after a Seek command is issued until the Sense Interrupt Status command is issued (following reception of the seek complete interrupt).

The FDC does not perform implied seeks. Before issuing data read or write commands, the read/write head must be positioned over the correct cylinder. If the head is not positioned correctly, a cylinder address error is generated.

After issuing a step pulse, the 8272 resumes drive status polling. For correct stepper operation in this mode, the stepper motor must be constantly enabled. (Most drives provide a jumper to permit the stepper motor to be constantly enabled.)

9. Step Rate

The 8272 can emit a step pulse that is one millisecond faster than the rate programmed by the SRT parameter in the Specify command. This action may cause subsequent sector not found errors. The step rate time should be programmed to be 1 ms longer than the step rate time required by the drive.

10. Cable Length

A cable length of less than 10 feet is recommended for drive interfacing.

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11. Scan Commands

The current 8272 has several problems when using the scan commands. These commands should not be used at this time.

12. Interrupts

When the processor receives an interrupt from the FDC, the FDC may be reporting one of two distinct events:

- a) The beginning of the result phase of a previously requested read, write, or scan command.
- b) An asynchronous event such as a seek/recalibrate completion, an attention, an abnormal command termination, or an invalid command.

These two cases are distinguished by the FDC busy flag (bit 4) in the Main Status Register. If the FDC busy flag is high, the interrupt is of type (a). If the FDC busy flag is low, the interrupt was caused by an asynchronous event (b).

A single interrupt from the FDC may signal more than one of the above events. After receiving an interrupt, the processor must continue to issue Sense Interrupt Status commands (and service the resulting conditions) until an invalid command code is received. In this manner, all "hidden" interrupts are ferreted out and serviced.

13. Skip Flag (SK)

The skip flag is used during the execution of Read Data, Read Deleted Data, Read Track, and various Scan commands. This flag permits the FDC to skip unwanted sectors on a disk track.

When performing a Read Data, Read Track, or Scan command, a high SK flag indicates that the FDC is to skip over (not transfer) any sector containing a deleted data address mark. A low SK flag indicates that the FDC is to terminate the command (after reading all the data in the sector) when a deleted data address mark is encountered.

When performing a Read Deleted Data command, a high SK flag indicates that sectors containing normal data address marks are to be skipped. Note that this is just the opposite situation from that described in the last paragraph. When a data address mark is encountered during a Read Deleted Data command (and the SK flag is low), the FDC terminates the command after reading all the data in the sector.

14. Bad Track Maintenance

The 8272 does not internally maintain bad track information. The maintenance of this information must be performed by system software. As an example of typical bad track operation, assume that a media test determines that track 31 and track 66 of a given floppy disk are bad. When the disk is formatted for use, the system software formats physical track 0 as logical cylinder 0 (C=0) in the command phase parameters), physical track 1 as logical track 1 (C=1), and so on, until physical track 30 is formatted as logical cylinder 30 (C=30). Physical track 31 is bad and should be formatted as logical cylinder FF (indicating a bad track). Next. physical track 32 is formatted as logical cylinder 31, and so on, until physical track 67 is formatted as logical cylinder 64. Next, bad physical track 66 is formatted as logical cylinder FF (another bad track marker), and physical track 67 is formatted as logical cylinder 65. This formatting continues until the last physical track (77) is formatted as logical cylinder 75. Normally, after this formatting is complete, the bad track information is stored in a prespecified area on the floppy disk (typically in a sector on track 0) so that the system will be able to recreate the bad track information when the disk is removed from the drive and reinserted at some later

To illustrate how the system software performs a transfer operation disk with bad tracks, assume that the disk drive head is positioned at track 0 and the disk described above is loaded into the drive. If a command to read track 36 is issued by an application program, the system software translates this read command into a seek to physical track 37 (since there is one bad track between 0 and 36, namely 31) followed by a read of logical cylinder 36. Thus, the cylinder parameter C is set to 37 for the Seek command and 36 for the Read Sector command.

15. Head Load versus Head Settle Times

The 8272 does not permit separate specification of the head load time and the head settle time. When the Specify command is issued for a given disk drive, the proper value for the HLT parameter is the maximum of the head load time and the head settle time.

APPENDIX

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Power Distribution

Power distribution						
Part	Ref Desig	+5	GND	+12	-12	
8088	A2	40	1,20			
8224	I6	9,16	8			
8237-2	A6	31	20]		
8251A	A9,B9,C9	26	4	1		
8253-5	A10	24	12	l		
8259A	B10	28	14		1	
8272	D10	40	20	}	1	
8284	. A1 ,	18	9			
8286	B6,F4	20	10			
2114	F1,F2,G1,G2,H1,H2,I1,I2	18	9			
2732A	D1,D2	- 24	12			
74LS00	E1	-14	7			
74LS04	B2,E6,E8,F8	14	1 7		٠ ,	
74LS27	E2,E5	14	1 7			
74LS32	B1	14	7			
74LS74	A4,G5,H6	14	7			
74LS138	F3	16	8			
74LS139	E10	16	8			
74LS153	13	16	8			
74LS157	F6	16	8			
74LS164	F5	14	. 7			
74LS173	G3	16	8	ł		
74LS175	G4	16	8			
74LS240	G10	20	10			
- 74LS257	D3	16	8		1	
74LS367	C3,E9	16	8			
74LS373	B4,C4,D4,C6	20	10 .			
74LS393	I5,F7	14	. 7			
74S08	E4	14	7			
74S138	D6,E3	16	8			
7414	Н7	14	7			
7438	H8,H9,H10	14	7			
1488	Н3		7 .	14	1.	
1489	H4 · _	14	7	, i		
96LS02	G7	16	8			
96LS02	G6 '			16	-8	
LM358	Н5		,	8	4	

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